

CS7800: Advanced Algorithms

Lecture 23: Regret minimization

- Jon's favorite algorithm
- Application to zero-sum games

Jonathan Ullman

12-9-2022

What do these problems have to do with each other?

- ① Finding equilibria in zero-sum games
- ② Solving linear programs
- ③ Minimizing a convex function $l(\theta)$
- ④ Boosting weak learning algorithms to better ensembles
- ⑤ Approximating SETCOVER
- ⑥ Generating privacy preserving synthetic data
- ⑦ Proving Chernoff bounds
- ⑧ Training many models without overfitting



No-Regret Learning

- Set of actions $\{1, 2, \dots, n\}$ you can take each day

- Play for T days

- For $t=1, 2, \dots, T$:

- You choose $P^t = (P_1^t, \dots, P_n^t)$ $\left[\sum_{i=1}^n P_i^t = 1, P_i^t \geq 0 \right]$

- You observe losses $l^t = (l_1^t, \dots, l_n^t)$ $[l_i^t \in [0, 1]]$

- You suffer loss $\sum_{i=1}^n l_i^t P_i^t \in [0, 1]$

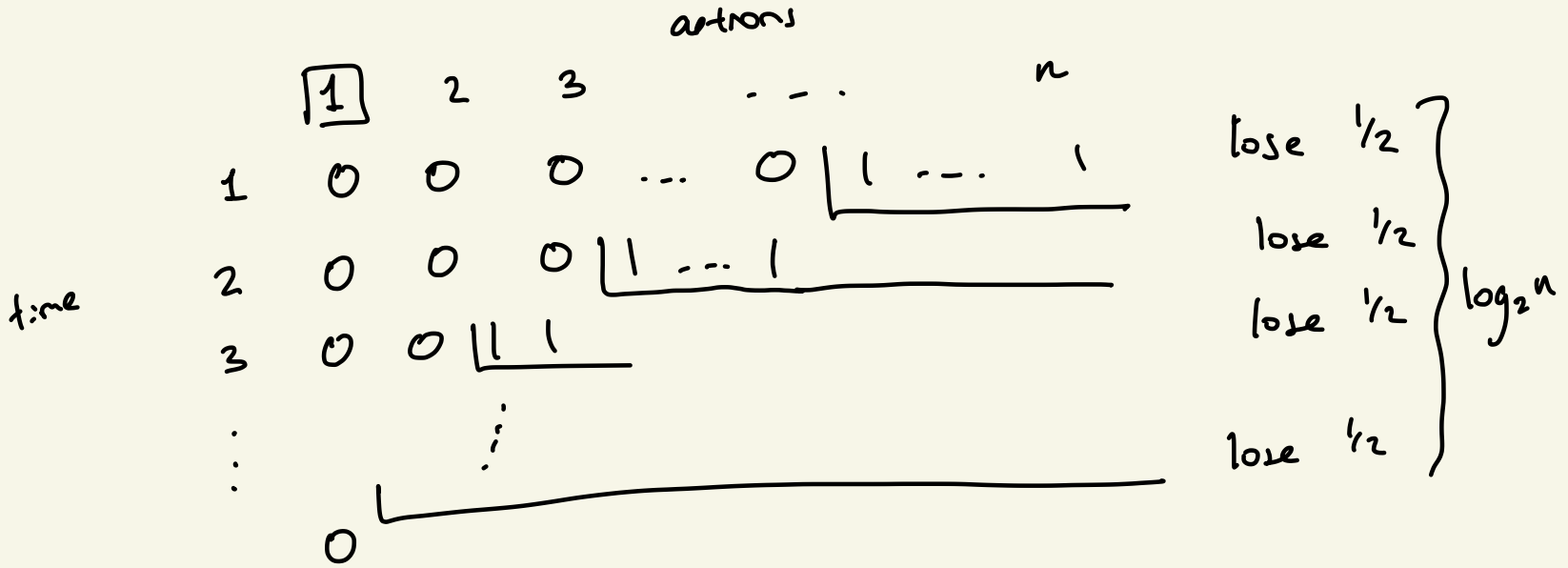
Regret

$$R = \underbrace{\sum_{t=1}^T \sum_{i=1}^n P_i^t l_i^t}_{L = \sum_{t=1}^T L^t} - \min_{j \in [n]} \underbrace{\sum_{t=1}^T l_j^t}_{L^*}$$

How to play when one action has no loss

- Initialize $w_i^1 = 1$ for $i=1, \dots, n$
 $W^1 = \sum_{i=1}^n w_i^1 = n$
- For $t=1, 2, \dots, T$:
 - Play strategy $P^t = \frac{w^t}{W^t}$
 - Receive losses $l^t \in \{0, 1\}^n$ and suffer loss $\underbrace{\sum_{i=1}^n p_i^t l_i^t}_{= L^t}$
 - Update: set $w_i^{t+1} = \begin{cases} 0 & \text{if } l_i^t = 1 \\ w_i^t & \text{if } l_i^t = 0 \end{cases}$

Worst Example



$$\frac{1}{n} + \frac{1}{n-1} + \frac{1}{n-2} + \dots + \frac{1}{2}$$

Analysis

$$\textcircled{1} W^1 = n \quad \textcircled{2} W^T \geq 1$$

$$\textcircled{3} \frac{W^t}{W^{t-1}} = (1 - L^{t-1})$$

$$1 \leq W^T = n \cdot \prod_{t=1}^T \frac{W^t}{W^{t-1}} = n \cdot \prod_{t=1}^T (1 - L^{t-1})$$

$$0 \leq \ln(n) + \sum_{t=1}^T \ln(1 - L^t) \leq \ln(n) - \sum_{t=1}^T L^t$$

$$\sum_{t=1}^T L^t \leq \ln(n)$$

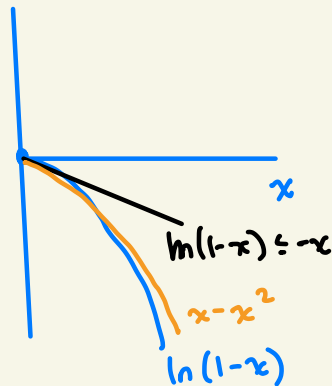
- Initialize $w_i^1 = 1$ for $i=1, \dots, n$
 $W^1 = \sum_{i=1}^n w_i^1 = n$

- For $t=1, 2, \dots, T$:

- Play strategy $P^t = \frac{w^t}{W^t}$

- Receive losses $l^t \in [0, 1]^n$ and suffer loss $\sum_{i=1}^n p_i^t l_i^t = L^t$

- Update: set $w_i^{t+1} = \begin{cases} 0 & \text{if } l_i^t = 1 \\ w_i^t & \text{if } l_i^t = 0 \end{cases}$



Multiplicative Weights

- Initialize $w_i^1 = 1$ for $i=1, \dots, n$
 $W^1 = \sum_{i=1}^n w_i^1 = n$

- For $t=1, 2, \dots, T$:

- Play strategy $P^t = \frac{w^t}{W^t}$

- Receive losses l^t and suffer loss $\sum_{i=1}^n p_i^t l_i^t = L^t$

- Update $w_i^{t+1} = (1-\eta)^{l_i^t} \cdot w_i^t$

$$W^{t+1} = \sum_{i=1}^n w_i^{t+1}$$

multiplicative weights updates

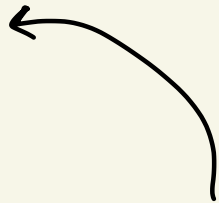
Analysis

① $W^1 = n$

② For every j : $W^T \geq W_j^T = \prod_{t=1}^T (1-\eta)^{l_j^t}$

③ $W^t = \sum_{i=1}^n w_i^t = \sum_{i=1}^n w_i^{t-1} \cdot (1-\eta)^{l_i^{t-1}} \leq \sum_{i=1}^n w_i^{t-1} \cdot (1-\eta L^{t-1})$

$$\frac{W^t}{W^{t-1}} \leq 1 - \eta L^{t-1}$$



- Initialize $w_i^1 = 1$ for $i=1, \dots, n$
 $W^1 = \sum_{i=1}^n w_i^1 = n$
- For $t=1, 2, \dots, T$:
 - Play strategy $P^t = \frac{w^t}{W^t}$
 - Receive losses l^t and suffer loss $\sum_{i=1}^n p_i^t l_i^t = L^t$
 - Update $w_i^{t+1} = (1-\eta)^{l_i^t} \cdot w_i^t$
 $W^{t+1} = \sum_{i=1}^n w_i^{t+1}$

$$\begin{aligned} &= \sum_{i=1}^n w_i^{t-1} - \eta w_i^{t-1} l_i^{t-1} \\ &= \sum_{i=1}^n w_i^{t-1} - W^{t-1} \cdot \eta P_i^{t-1} l_i^{t-1} \\ &= W^{t-1} - W^{t-1} \eta L^{t-1} \\ &= W^{t-1} (1 - \eta L^{t-1}) \end{aligned}$$

Analysis

① $W^1 = n$

② For every j : $W^T \geq \omega_j^T = \prod_{t=1}^T (1-\eta)^{l_j^t}$

③ $\frac{W^t}{W^{t-1}} \leq 1 - \eta L^{t-1}$

$$\prod_t (1-\eta)^{l_{j^*}^t} \leq \cancel{W^T} \leq n \cdot \prod_t (1-\eta L^{t-1})$$

$$L^* (1-\eta - \eta^2) \leq \sum_t l_{j^*}^t \ln(1-\eta) \leq \ln(n) + \sum_t \ln(1-\eta L^{t-1}) \leq \ln(n) - \eta L$$

$$-\eta L^* - \eta^2 L^* \leq \ln(n) - \eta L$$

$$\eta(L - L^*) \leq \ln(n) + \eta^2 L^*$$

$$\rightarrow L - L^* \leq \frac{\ln(n)}{\eta} + \eta L^* \leq \frac{\ln(n)}{\eta} + \eta T$$

$$\eta = \sqrt{\frac{\ln(n)}{T}} \Rightarrow L - L^* \leq 2\sqrt{T \ln(n)}$$

• Initialize $\omega_i^1 = 1$ for $i=1, \dots, n$
 $W^1 = \sum_{i=1}^n \omega_i^1 = n$

• For $t=1, 2, \dots, T$:

- Play strategy $P^t = \frac{W^t}{W^t}$

- Receive losses l^t and suffer loss $\sum_{i=1}^n p_i^t l_i^t = L^t$

- Update $\omega_i^{t+1} = (1-\eta)^{l_i^t} \cdot \omega_i^t$
 $W^{t+1} = \sum_{i=1}^n \omega_i^{t+1}$