

CS7800: Advanced Algorithms

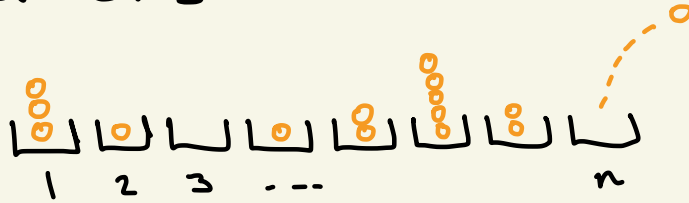
Lecture 17: Randomized Algorithms II

- Finish balls and bins
- String matching

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Balls and Bins



Throw m balls into n
bins independently

$$\omega = (1, 8, 19, 23, 6, 2)$$

ball 1 to bin 1
ball 2 to bin 8
ball 3 to bin 19

$$P(\omega_1, \omega_2, \dots, \omega_m) = \frac{1}{n^m}$$

Questions:

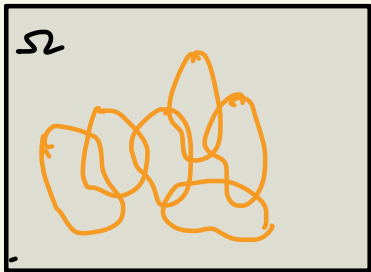
- ① How long until bin 1 gets a ball (in expectation)?
- ② How long until no bins are empty (in expectation)?
- ③ What is the most number of balls in any bin (in expectation)?

Maximum Load



want to understand $IP(M \gg t)$
or $IE(M)$

- $M(\omega) = \text{maximum load}$
- $M_j = \text{load on bin } j$ $M = \max \{M_1, M_2, \dots, M_n\}$
- $IP(M \gg t) = IP((M_1 \gg t) \vee (M_2 \gg t) \vee \dots \vee (M_n \gg t))$



"union bound"

$$\leq \sum_{j=1}^n IP(M_j \gg t) = n \cdot \underline{IP(M_1 \gg t)}$$

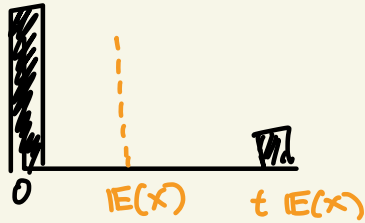
Want to find t
so that this is $\ll \frac{1}{n}$

Markov's Inequality

Let X be any random variable such that $X(\omega) \geq 0$

Thm (Markov): $P(X \geq t) \leq \frac{E(X)}{t}$

Pf.



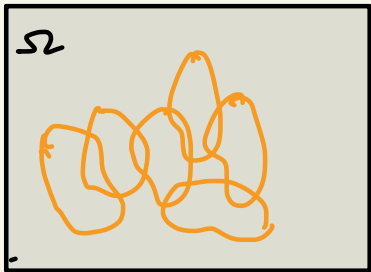
$$\begin{aligned} E(X) &= \sum_{i=0}^{\infty} i \cdot P(X=i) = \sum_{i=0}^{t-1} i \cdot P(X=i) + \sum_{i=t}^{\infty} i \cdot P(X=i) \\ &\geq 0 + \sum_{i=t}^{\infty} t \cdot P(X=i) \\ &= t \cdot P(X \geq t) \end{aligned}$$

Maximum Load



want to understand $IP(M \gg t)$
or $IE(M)$

- $M(\omega) = \text{maximum load}$
- $M_j = \text{load on bin } j$ $M = \max \{M_1, M_2, \dots, M_n\}$
- $IP(M \gg t) = IP((M_1 \gg t) \vee (M_2 \gg t) \vee \dots \vee (M_n \gg t))$



“union bound” \rightarrow

$$\leq \sum_{j=1}^n IP(M_j \gg t) = n \cdot IP(M_1 \gg t)$$

Want to find t so that this is $\ll \frac{1}{n}$

$$n \cdot IP(M_1 \gg t) \leq n \cdot \frac{IE(M_1)}{t}$$
$$= \frac{n \cdot (m/n)}{t} = \frac{m}{t}$$

Chebyshev's Inequality

Let X be any random variable, let $\mu = \mathbb{E}(X)$

Thm: $\mathbb{P}(|X - \mu| \geq t) \leq \frac{\mathbb{E}((X - \mu)^2)}{t^2} = \frac{\text{Var}(X)}{t^2}$

Pf: $\mathbb{P}(|X - \mu| \geq t) = \mathbb{P}((X - \mu)^2 \geq t^2) \leq \frac{\mathbb{E}((X - \mu)^2)}{t^2}$

↑
Markov

Applying Chebyshev to Balls and Bins

Let M_i be the load of bin i

$$\mathbb{E}(M_i) = \frac{m}{n} \quad (=1)$$

Let $M_{i,j} = \begin{cases} 1 & \text{if ball } j \text{ goes to bin } i \\ 0 & \text{otherwise} \end{cases}$

indicator
random
variable

$$\mathbb{E}((M_i - \mu)^2) = \mathbb{E}(M_i^2) - \mathbb{E}(M_i)^2 \leq \frac{m}{n} \quad (=1)$$

$$\mathbb{E}(M_i^2) = \mathbb{E}((M_{i,1} + \dots + M_{i,m})^2) \quad \mathbb{E}(M_i)^2 = \frac{m^2}{n^2} \quad (=1)$$

$$= \mathbb{E}\left(\sum_{i=1}^m M_{i,i}^2 + \sum_{i \neq j} M_{i,i} M_{i,j}\right)$$

$$= m \cdot \frac{1}{n} + m(m-1) \cdot \frac{1}{n^2}$$

$$\leq \frac{m}{n} + \frac{m^2}{n^2} \quad (=2)$$

$$\mathbb{E}(I) = \mathbb{P}(I)$$

Applying Chebyshev to Balls and Bins

Let M_i be the load of bin i

$$\mathbb{E}(M_i) = \frac{m}{n}$$

$m = n$

$$\left(= 1 \right)$$

Let $M_{i,j} = \begin{cases} 1 & \text{if ball } j \text{ goes to bin } i \\ 0 & \text{otherwise} \end{cases}$

indicator
random
variable

$$\mathbb{E}((M_i - \mu)^2) = \mathbb{E}(M_i^2) - \mathbb{E}(M_i)^2 \leq \frac{m}{n} \quad (\leq 1)$$

$$\mathbb{E}(I) = P(I)$$

$$\begin{aligned} P(M_i \geq t) &\leq n \cdot P(M_{i,j} \geq t) \leq n \cdot \frac{\mathbb{E}((M_{i,j} - \frac{m}{n})^2)}{t^2} \\ &\leq n \cdot \frac{m/n}{t^2} = \frac{m}{t^2} \end{aligned}$$

$$P(M_i \geq 10 \cdot \sqrt{m}) \leq \frac{1}{100} \quad (\text{Max load is at most } \approx \sqrt{m})$$

Chernoff Bounds

Let Z_1, \dots, Z_n be independent r.v.'s such that

$$Z_i = \begin{cases} 1 & \text{w.p. } p_i \\ 0 & \text{w.p. } 1-p_i \end{cases}$$

and $Z = Z_1 + \dots + Z_n$. Let $\mu = \mathbb{E}(Z) = p_1 + \dots + p_n$

Thm: $\mathbb{P}(Z > (1+\delta)\mu) \leq \left(\frac{e^\delta}{(1+\delta)^{(1+\delta)}} \right)^\mu$

Z has mean μ and variance $\sum_{i=1}^n p_i(1-p_i) \approx \mu$

$$Z \approx \mathcal{N}\left(\mu, \sum_i p_i(1-p_i)\right)$$

Chernoff Bound Proof

Let z_1, \dots, z_n be independent r.v.'s such that

$$z_i = \begin{cases} 1 & \text{w.p. } p_i \\ 0 & \text{w.p. } 1-p_i \end{cases}$$

and $z = z_1 + \dots + z_n$. Let $\mu = \mathbb{E}(z) = p_1 + \dots + p_n$

Thm: $\mathbb{P}(z > (1+\delta)\mu) \leq \left(\frac{e^\delta}{(1+\delta)^{1+\delta}}\right)^\mu$

$$\mathbb{P}(z > a\mu) = \mathbb{P}(e^{tz} > e^{ta\mu})$$

$$\leq e^{-ta\mu} \cdot \mathbb{E}(e^{tz})$$

$$= e^{-ta\mu} \cdot \prod_{i=1}^n \mathbb{E}(e^{tz_i}) = e^{-ta\mu} \cdot \prod_{i=1}^n (p_i e^t + 1 - p_i) = e^{-ta\mu} \cdot \prod_{i=1}^n (1 + p_i(e^t - 1))$$

$$\leq e^{-ta\mu} \cdot \prod_{i=1}^n e^{p_i(e^t - 1)} = e^{-ta\mu} \cdot e^{(e^t - 1) \cdot \sum_{i=1}^n p_i} = e^{-ta\mu} \cdot e^{(e^t - 1)\mu}$$

$$= e^{(e^t - 1)\mu - ta\mu}$$

$$(a = (1+\delta))$$

$$= \left(e^{e^t - 1 - ta}\right)^\mu \leftarrow$$

Set $a = 1+\delta$, plug in the right value of t

Applying Chernoff to Balls and Bins

M_i = load on bin i

$$M_i = M_{i,1} + \dots + M_{i,m}$$

where $M_{i,j}$ is an indicator

$$\mathbb{E}(M_{i,j}) = \frac{1}{n}$$

Let Z_1, \dots, Z_n be independent r.v.'s such that

$$Z_i = \begin{cases} 1 & \text{w.p. } p_i \\ 0 & \text{w.p. } 1-p_i \end{cases}$$

and $Z = Z_1 + \dots + Z_n$. Let $\mu = \mathbb{E}(Z) = p_1 + \dots$

Thm: $\mathbb{P}(Z > (1+\delta)\mu) \leq \left(\frac{e^\delta}{(1+\delta)^{1+\delta}}\right)^\mu$

$$\mathbb{P}(M_i > (1+\delta)\frac{m}{n}) \leq n \cdot \mathbb{P}(M_{i,j} > (1+\delta)\frac{m}{n}) \leq n \cdot \left(\frac{e^\delta}{(1+\delta)^{1+\delta}}\right)^{\frac{m}{n}}$$

m/n "large"

$$\mathbb{P}(M_i > 3 \cdot \frac{m}{n}) \leq n \cdot (.9)^{\frac{m}{n}} \stackrel{\text{Assume } \frac{m}{n} \geq 10 \log n}{\leq} n \cdot \frac{1}{n^2} = \frac{1}{n}$$

$\frac{m}{n}$ is "small" e.g. $n=n$

if $\delta \approx \log n$

$$\mathbb{P}(M_i > \log n) \leq n \cdot \left(\frac{e^{\log n}}{(\log n)^{1+\log n}}\right) \leq \frac{1}{n}$$